

# Equilibration and Unitary $k$ - Designs

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UCL

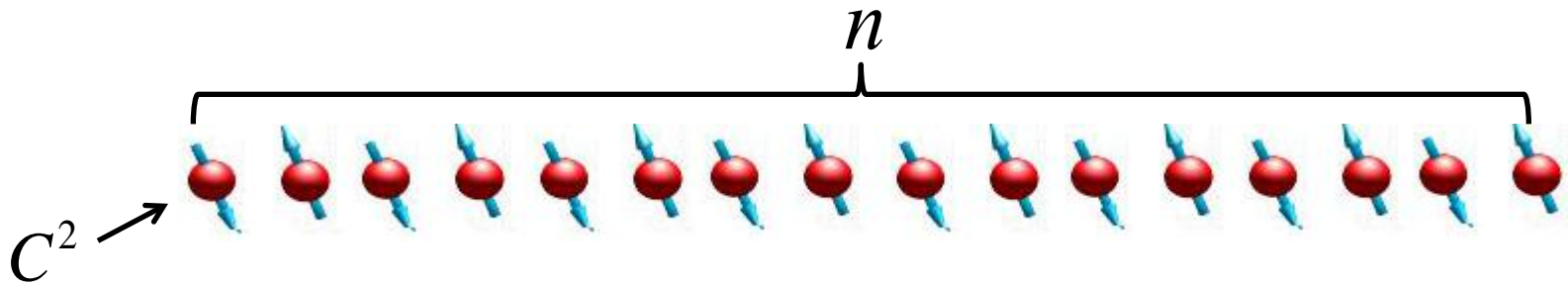
Joint work with

Aram Harrow and Michal Horodecki

arXiv:1208.0692

IMS, September 2013

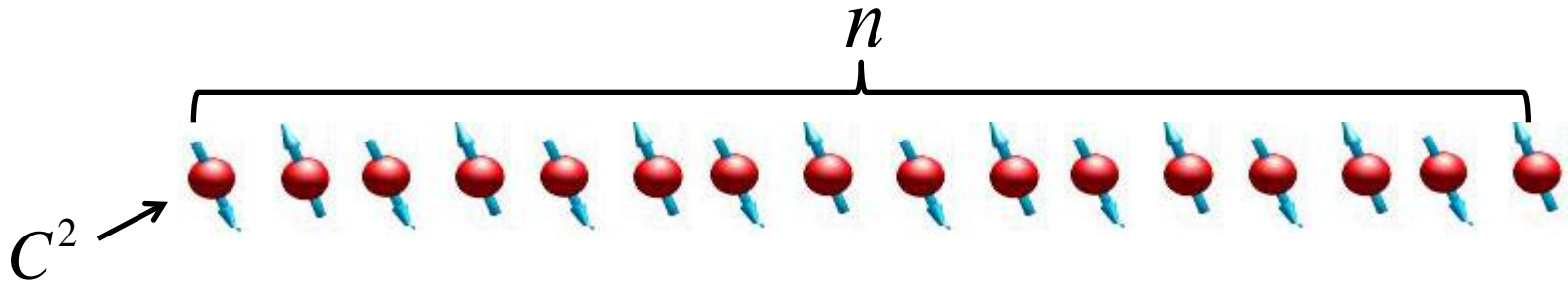
# Dynamical Equilibration



State at time  $t$ :  $|\psi_t\rangle = U(t)|0\rangle^{\otimes n}$

$$U(t) = \mathcal{T} e^{i \int_0^t H(t') dt'}$$

# Dynamical Equilibration



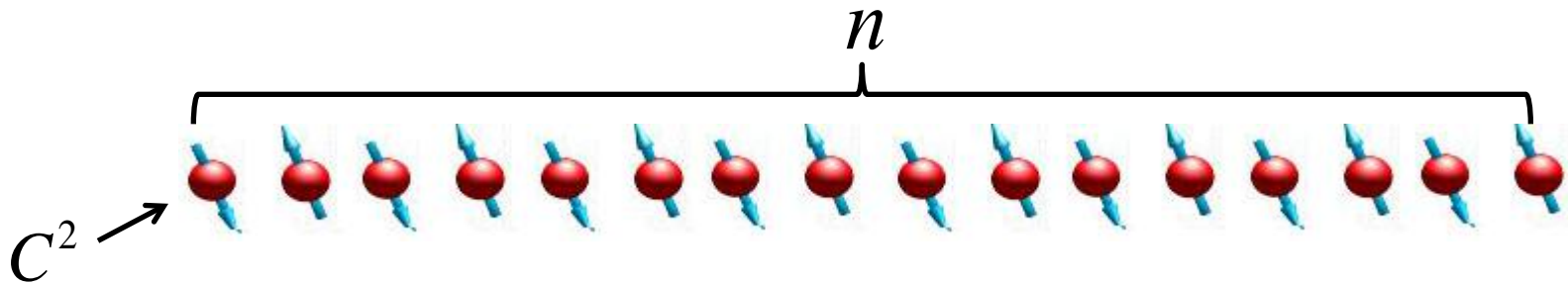
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I.e. for most  $t$   $|\psi(t)\rangle\langle\psi(t)| \approx \rho_0$  ?

# Dynamical Equilibration



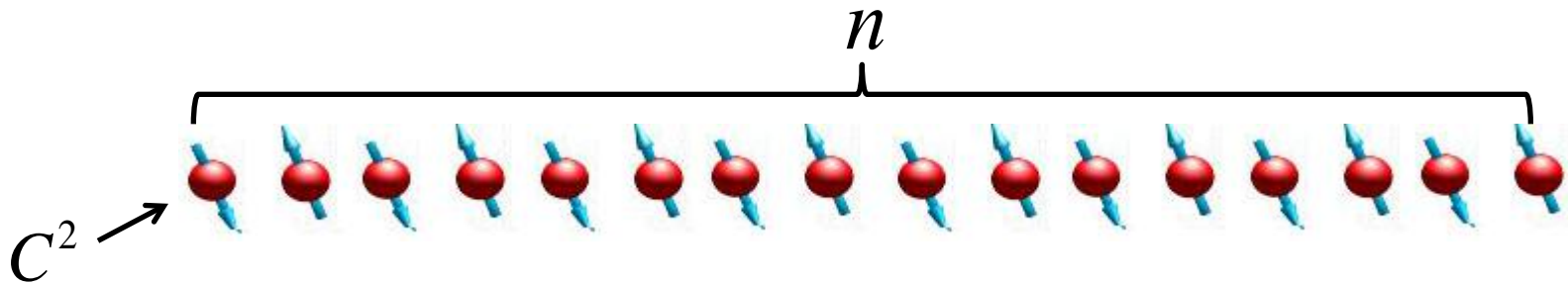
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$$U(t) = \mathcal{T} e^{i \int_0^t H(t') dt'}$$

Will  $|\psi(t)\rangle$  equilibrate?

I.e. for most  $t$   $|\psi(t)\rangle\langle\psi(t)| \approx \rho_0$  ? **NO!**

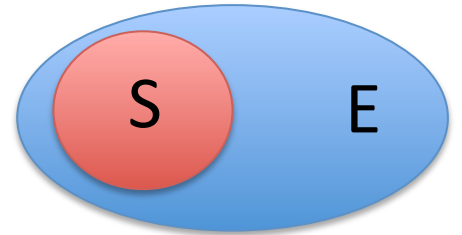
# Dynamical Equilibration



How about relative to particular kind of measurements?

- “macroscopic” measurements (von Neumann ‘29)
- local measurements
- local measurements relative to an external observer (Lidia’s talk)
- Low-complexity measurements (i.e. measurements that require time much less than  $t$ )

# Equilibration is generic



(Linden, Popescu, Short, Winter '08)

Almost any Hamiltonian  $H$  with equilibrate:

$$\mathbb{E}_{t \geq 0} \|\rho_S(t) - \overline{\rho_S}\|_1 \approx 0$$

with  $\rho_S(t) = \text{tr}_{\setminus S}(e^{itH} |0^n\rangle\langle 0^n| e^{-itH})$

and  $\overline{\rho_S} = \lim_{T \rightarrow \infty} \mathbb{E}_{t \leq T} \rho_S(t)$

# Time Scale of Equilibration

The previous approach only gives bounds exponentially small in the number of particles

**Can we prove *fast* equilibration is generic?**

For particular cases better bounds are known

E.g. (Cramer et al '08), (Banuls, Cirac, Hastings '10), ...

This talk:

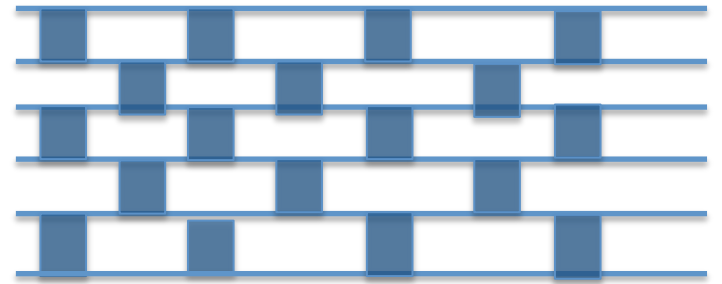
**Generic local dynamics leads to rapid equilibration**

Caveat: time-dependent Hamiltonians...



# Parallel Random Quantum Circuits

**Parallel Local Random Circuit:** in each step  $n/2$  independent Haar two-qubit gates are applied to either  $((1, 2), (3, 4), \dots, (n-1, n))$  or  $((2, 3), (4, 5), \dots, (n-2, n-1))$



Discrete version of  $U(t) = \mathcal{T} e^{i \int_0^t H(t') dt'}$

with random  $H(t) = H_{12}(t) + H_{23}(t) + \dots + H_{n-1}(t)$

# Equilibration for Random Circuits

**Thm** Let  $RC_t := \{U : U = U_1 \dots U_t\}$  be the set of all circuits of length  $t$

1. For every region  $S$ ,  $\mathbb{E}_{U \in RC_t} \|\rho_S(t) - \tau\|_1 \leq \varepsilon$

for  $t \geq O(\log(1/\varepsilon)\text{size}(S))$

with  $\rho_S(U) := \text{tr}_{\setminus S}(U|0\rangle\langle 0|^{\otimes n}U^\dagger)$



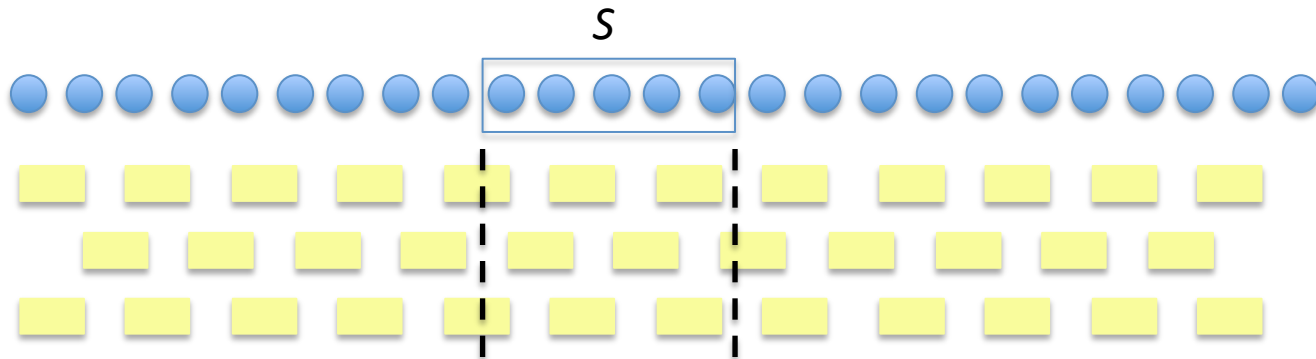
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The result matches the speed-of-sound propagation bound

# Equilibration for Random Circuits

**Thm** Let  $RC_t := \{ U : U = U_1 \dots U_t \}$  be the set of all circuits of length  $t$

2. Let  $M_k := \{ M : 0 \leq M \leq \text{id}, M \text{ has gate complexity } k \}$ .

For every  $t \geq O(k^6 \log(1/\varepsilon))$

$$\mathbb{E}_{U \in RC_t} \max_{N \in M_k} \left| \langle 0^n | U^\dagger N U | 0^n \rangle - \text{tr}(N)/2^n \right| \leq \varepsilon$$

Equilibration for arbitrary measurements of *low* complexity

Fails for  $t = k$ , as one can undo the evolution  $U$

# Warm-up: Equilibration for Haar Random Unitaries

(Page '93)

Haar measure

Let  $|\psi\rangle = U|0\rangle^{\otimes n}$  and  $\rho_S(U) = \text{tr}_{S^c}(U|0\rangle\langle 0|^{\otimes n}U^\dagger)$

We have  $\|\rho_S(U) - \tau\|_1 \leq \dim(S)^{1/2} \|\rho_S(U) - \tau\|_2$

But  $\mathbb{E}_U \|\rho_S(U) - \tau\|_2^2 = \frac{\dim(S)}{\dim(S^c)}$  only second moments needed

So  $\|\rho_S(U) - \tau\|_1 \leq \varepsilon$  for  $\log(\dim(S^c)) \geq 2 \log(\dim(S)/\varepsilon)$



# Unitary $k$ -designs

**Def.** An ensemble of unitaries  $\{\mu(dU), U\}$  in  $\mathbf{U}(d)$  is an  $\varepsilon$ -approximate unitary  $k$ -design if for every monomial

$$M = U_{p_1, q_1} \dots U_{p_k, q_k} U_{r_1, s_1}^* \dots U_{r_k, s_k}^*$$

$$|E_{\mu}(M(U)) - E_{\text{Haar}}(M(U))| \leq \varepsilon$$

First  $k$  moments are close to the Haar measure

Equivalent to ( $\approx$ )

$$\left\| \int \mu(dU) U^{\otimes k} \otimes \bar{U}^{\otimes k} - \int \mu_{\text{Haar}}(dU) U^{\otimes k} \otimes \bar{U}^{\otimes k} \right\|_{\infty} \leq \varepsilon$$

# Unitary $k$ -designs

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$$|E_{\mu}(M(U)) - E_{\text{Haar}}(M(U))| \leq \varepsilon$$

Natural quantum generalization of  $k$ -wise independent distributions

Many applications in quantum information theory: encoding for quantum communication (2-design), generic speed-ups (3-design), efficient tomography (4-design), ...

# Unitary $k$ -designs and equilibration

From  $\delta$ -approx 2-design

Let  $|\psi\rangle = U|0\rangle^{\otimes n}$  and  $\rho_S(U) = \text{tr}_{S^c}(U|0\rangle\langle 0|^{\otimes n}U^\dagger)$

We have  $\|\rho_S(U) - \tau\|_1 \leq \dim(S)^{1/2} \|\rho_S(U) - \tau\|_2$

But  $\mathbb{E}_U \|\rho_S(U) - \tau\|_2^2 = \frac{\dim(S)}{\dim(S^c)} + \delta$

So  $\|\rho_S(U) - \tau\|_1 \leq \sqrt{\frac{\dim(S)}{\dim(S^c)} + \delta} \dim(S)$



# Equilibration for Random Circuits

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2. Let  $M_k := \{ M : 0 \leq M \leq \text{id}, M \text{ has gate complexity } k \}$ .

For every  $t \geq O(k^6 \log(1/\varepsilon))$

$$\mathbb{E}_{U \in RC_t} \max_{N \in M_k} \left| \langle 0^n | U^\dagger N U | 0^n \rangle - \text{tr}(N)/2^n \right| \leq \varepsilon$$

Proof follows by looking at higher moments to obtain a good concentration bound on  $\langle 0^n | U^\dagger N U | 0^n \rangle$  and union bound over set  $M_k$ .

Requires approximate **poly(k)**-design

# Unitary $k$ -designs

## Previous work:

(DiVincenzo, Leung, Terhal '02) Clifford group is an *exact* 2-design

(Dankert et al '06) Efficient construction of 2-design

(Ambainis and Emerson '07) Efficient construction of *state*  
poly( $n$ )-design

(Harrow and Low '08) Efficient construction of ( $n/\log(n)$ )-design

(Abeyesinghe '06) 2-designs are enough for *decoupling*

(Low '09) Other applications of  $t$ -design (mostly 2-designs)  
replacing Haar unitaries

# Random Quantum Circuits vs Unitary Designs

## Previous work:

(Oliveira, Dalhsten, Plenio '07)  $O(n^3)$  random circuits are 2-designs

(Harrow, Low '08)  $O(n^2)$  random Circuits are 2-designs for every universal gate set

(Arnaud, Braun '08) numerical evidence that  $O(n \log(n))$  random circuits are unitary  $t$ -design

(Znidaric '08) connection with spectral gap of a mean-field Hamiltonian for 2-designs

(Brown, Viola '09) connection with spectral gap of Hamiltonian for  $t$ -designs

(B., Horodecki '10)  $O(n^2)$  local random circuits are 3-designs

# Random Quantum Circuits as $k$ -designs?

Conjecture Random Circuits of size  $\text{poly}(n, \log(1/\epsilon))$  are an  $\epsilon$ -approximate unitary  $\text{poly}(n)$ -design

# Random Quantum Circuits as $k$ -designs?

Thm 1 Local Random Circuits of size  $O(nk^4 \log(1/\epsilon))$  are an  $\epsilon$ -approximate unitary  $k$ -design

Thm 2 Parallel Local Random Circuits of size  $O(k^4 \log(1/\epsilon))$  are an  $\epsilon$ -approximate unitary  $k$ -design

# Equilibration for Random Circuits

**Thm** Let  $RC_t := \{U : U = U_1 \dots U_t\}$  be the set of all circuits of length  $t$

1. For every region  $S$ ,  $\mathbb{E}_{U \in RC_t} \|\rho_S(t) - \tau\|_1 \leq \varepsilon$

for  $t \geq O(\log(1/\varepsilon)\text{size}(S))$

with  $\rho_S(U) := \text{tr}_{\setminus S}(U|0\rangle\langle 0|^{\otimes n}U^\dagger)$

Proof follows from the calculation for a 2-design from before

# Equilibration for Random Circuits

**Thm** Let  $RC_t := \{ U : U = U_1 \dots U_t \}$  be the set of all circuits of length  $t$

2. Let  $M_k := \{ M : 0 \leq M \leq \text{id}, M \text{ has gate complexity } k \}$ .

For every  $t \geq O(k^6 \log(1/\varepsilon))$

$$\mathbb{E}_{U \in RC_t} \max_{N \in M_k} \left| \langle 0^n | U^\dagger N U | 0^n \rangle - \text{tr}(N)/2^n \right| \leq \varepsilon$$

Proof follows by looking at higher moments to obtain a good concentration bound on  $\langle 0^n | U^\dagger N U | 0^n \rangle$ , for fixed  $N$ , and take the union bound over the set  $M_k$ .

Requires approximate  $\text{poly}(k)$ -design

# Outline Proof of Main Result

1. Mapping the problem to bounding spectral gap of a **Local Hamiltonian**
2. Technique for **bounding spectral gap** (Nachtergaele '94) + representation theory  
(reduces the problem to obtaining an exponentially small lower bound on the spectral gap)
3. **Path Coupling** applied to the unitary group  
(prove convergence of the random walk in exponential time)
4. Use **detectability Lemma** (Arad et al '10) to go from local random circuits to *parallel* local random circuits

# Relating to Spectral Gap

$\mu_n$ : measure on  $U(2^n)$  induced by one step of the local random circuit model

$(\mu_n)^{*k}$ :  $k$ -fold convolution of  $\mu_n$  (measure induced by  $k$  steps of the local random circuit model)

By eigendecomposition

$$\int \mu_n(dU) U^{\otimes t,t} = \int \mu_{\text{Haar}}(dU) U^{\otimes t,t} + \sum_l \lambda_l P_l, \quad |\lambda_l| < 1$$

$$\text{so } \left\| \int \mu_n^{*t}(dU) U^{\otimes t,t} - \int \mu_{\text{Haar}}(dU) U^{\otimes t,t} \right\|_{\infty} = \lambda_2^t$$

# Relating to Spectral Gap

$\mu_n$ : measure of circuit model

It suffices to prove an upper bound on  $\lambda_2$  of the form  $1 - \Omega(t^{-4}n^{-1})$  since  $(1 - \Omega(t^{-4}n^{-2}))^k \leq \epsilon$  for  $k = O(nt^4 \log(1/\epsilon))$

$(\mu_n)^{*k}$ :  $k$ -fold local random

By eigendecomposition

$$\int \mu_n(dU) U^{\otimes t,t} = \int \mu_{\text{Haar}}(dU) U^{\otimes t,t} + \sum_l \lambda_l P_l, \quad |\lambda_l| < 1$$

$$\text{so } \left\| \int \mu_n^{*t}(dU) U^{\otimes t,t} - \int \mu_{\text{Haar}}(dU) U^{\otimes t,t} \right\|_{\infty} = \lambda_2^t$$

# Relating to Spectral Gap

But 
$$\mu_n = \frac{1}{n} \sum_{i=1}^n \mu_{Haar}(i, i+1)$$

So 
$$\lambda_2 \left( \int \mu_n(dU) U^{\otimes t} \otimes \bar{U}^{\otimes t} \right) = 1 - \frac{\Delta(H_{n,t})}{n}$$

with 
$$H_{n,t} := \sum_{i=1}^n h_{i,i+1} \quad h_{i,i+1} := I - \int_{U(4)} U_{i,i+1}^{\otimes t} \otimes \bar{U}_{i,i+1}^{\otimes t} \mu_H(dU)$$

and  $\Delta(H_{n,t})$  the spectral gap of the local Hamiltonian  $H_{n,t}$



# Relating to Spectral Gap

But 
$$\mu_n = \frac{1}{n} \sum_{i=1}^n \mu_{Haar}(i, i+1)$$

Want to lower bound spectral gap by  $O(t^{-4})$

So 
$$\lambda_2 \left( \int \mu_n(dU) U^{\otimes t} \otimes \bar{U}^{\otimes t} \right) = 1 - \frac{\Delta(H_{n,t})}{n}$$

with 
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and  $\Delta(H_{n,t})$  the spectral gap of the local Hamiltonian  $H_{n,t}$



# Structure of $H_{n,t}$

$$H_{n,t} := \sum_{i=1}^n h_{i,i+1} \quad h_{i,i+1} := I - \int_{U(4)} U_{i,i+1}^{\otimes t} \otimes \bar{U}_{i,i+1}^{\otimes t} \mu_H(dU)$$

i.  $H_{n,t} \in ((\mathbb{C}^2)^{\otimes 2t})^{\otimes n}$ , frustration-free with min eigenvalue 0

ii.  $\int \mu_{\text{Haar}}(dU) U^{\otimes t} \otimes \bar{U}^{\otimes t}$  projects onto 0 eigenspace  $G_{n,t}$ :

$$G_{n,t} := \text{span} \left\{ |\psi_\pi\rangle^{\otimes n}, |\psi_\pi\rangle := (I \otimes V(\pi)) |\Phi(2^t)\rangle : \pi \in S_t \right\}$$

$$\begin{aligned} & V(\pi) |a_1, \dots, a_t\rangle \\ &= |a_{\pi^{-1}(1)}, \dots, a_{\pi^{-1}(t)}\rangle \end{aligned}$$

$$|\Phi(d)\rangle = d^{-1/2} \sum_{i=1}^d |i, i\rangle$$

# Approximate Orthogonality

$|\psi_\pi\rangle$  are non-orthogonal, but

$$\sum_{\pi \in S_t} |\langle \psi_\sigma | \psi_\pi \rangle|^n \leq 1 + \frac{2t^2}{2^n}, \quad \left\| \sum_{\pi \in S_t} (|\psi_\pi\rangle\langle\psi_\pi|)^{\otimes n} - G_{n,t} \right\|_\infty \leq \frac{2t^2}{d^n}$$

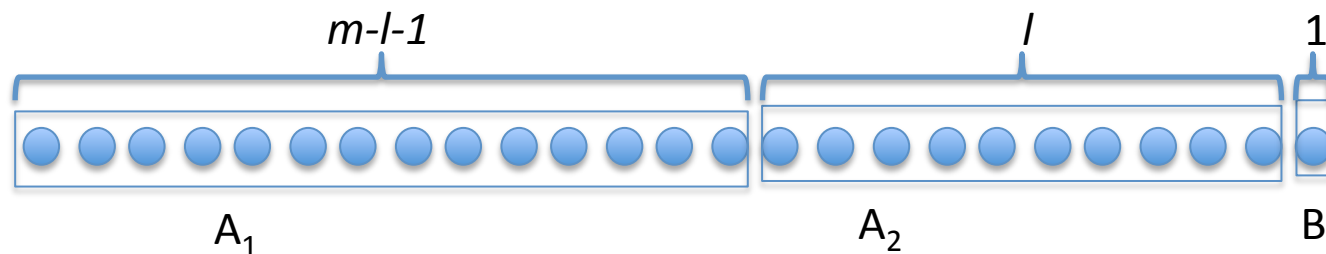
Proof by basic representation theory, in fact only uses

$$P_{sym} = \frac{1}{t!} \sum_{\pi \in S_t} V(\pi)$$

# Lower Bounding $\Delta(H_{n,t})$

Lemma:  $\Delta(H_{n,t}) \geq \frac{\Delta(H_{\log(t),t})}{8 \log(t)}$

Follows from structure of  $H_{n,t}$ , approx. orthogonality, and



(Nachtergaele '96) Suppose there exists  $l$  and  $\varepsilon_l < l^{-1/2}$  s.t.

$$\left\| (I_{A_1} \otimes G_{A_2 B})(G_{A_1 A_2} \otimes I_B) - G_{A_1 A_2 B} \right\|_{\infty} \leq \varepsilon_l$$

Then:  $\Delta(H_{[1,n]}) \geq \Delta(H_{[1,l]}) \left( \frac{(1 - \varepsilon_l \sqrt{l})}{l-1} \right)$

# Lower Bounding $\Delta(H_{n,t})$

Lemma:  $\Delta(H_{n,t}) \geq \frac{\Delta(H_{\log(t),t})}{8 \log(t)}$

Follow

Want to lower bound by  $O(t^{-4})$ , an exponential small bound in the size of the chain (i.e. in  $\log(t)$ )

$A_1$

$A_2$

$B$

(Nachtergaele '96) Suppose there exists  $l$  and  $\varepsilon_l < l^{-1/2}$  s.t.

$$\left\| (I_{A_1} \otimes G_{A_2 B})(G_{A_1 A_2} \otimes I_B) - G_{A_1 A_2 B} \right\|_{\infty} \leq \varepsilon_l$$

Then:  $\Delta(H_{[1,n]}) \geq \Delta(H_{[1,l]}) \left( \frac{(1 - \varepsilon_l \sqrt{l})}{l - 1} \right)$

# Exponentially Small Bound to Spectral Gap

Follows from two relations:

$$1. \left(1 - \frac{\Delta(H_{n,t})}{n}\right)^k \leq 2tW\left(\left(\mu_n\right)^{*k}, \mu_{Haar}\right)$$

Wasserstein distance:

$$W(\nu_1, \nu_2) := \sup\left\{\int f(u)\nu_1(du) - \int f(u)\nu_2(du) : f \text{ is } 1\text{-Lipschitz}\right\}$$

$$2. W\left(\left(\mu_n\right)^{*{(n-1)k}}, \mu_{Haar}\right) \leq 2^{n/2} (1 - 2^{-5n})^{\frac{k}{n-1}}$$

# Bounding Convergence with Path Coupling

**Key result** to 2<sup>nd</sup> relation: Extension to the unitary group of Bubley and Dyer **path coupling**

Let  $W_p(\nu_1, \nu_2) := \inf \left\{ E[d(X, Y)^p]^{1/p} : (X, Y) \text{ couples } (\nu_1, \nu_2) \right\}$

(Oliveira '07) Let  $\nu$  be a measure in  $U(d)$  s.t.

$$\lim_{\varepsilon \rightarrow 0} \sup_{U_1, U_2 \in U(d)} \left\{ \frac{W_2(\nu * \delta_{U_1}, \nu * \delta_{U_2})}{\|U_1 - U_2\|_2} : \|U_1 - U_2\|_2 \leq \varepsilon \right\} \leq \eta$$

Then  $W_2(\nu * \nu_1, \nu * \nu_2) \leq \eta W_2(\nu_1, \nu_2)$

# Bounding Convergence with Path Coupling

Must consider coupling in  $n$  steps of the walk to get non trivial contraction (see paper for details)

(Oliveira '07) Let  $\nu$  be a measure in  $U(d)$  s.t.

$$\lim_{\varepsilon \rightarrow 0} \sup_{U_1, U_2 \in U(d)} \left\{ \frac{W_2(\nu * \delta_{U_1}, \nu * \delta_{U_2})}{\|U_1 - U_2\|_2} : \|U_1 - U_2\|_2 \leq \varepsilon \right\} \leq \eta$$

Then  $W_2(\nu * \nu_1, \nu * \nu_2) \leq \eta W_2(\nu_1, \nu_2)$

# Time-Independent Models

Toy model for equilibration: Let  $H_{SE} = UDU^T$ , with  $U$  taken from the Haar measure in  $U(|S| |E|)$  and  $D := \text{diag}(E_1, E_2, \dots)$ .

$$\frac{1}{T} \int_{U(|S||E|)} \text{tr} \left( (\rho_S(t) - \rho_{\text{equi}})^2 \right) \mu_{\text{Haar}}(dU) \leq O \left( \frac{\left( \sum_k e^{i2tE_k} \right)^2}{|S||E|^2} + \dots \right)$$

(B., Ciwiklinski et al '11, Masanes et al '11, Vinayak, Znidaric '11)

Time of equilibration: Average energy gap:  $\frac{1}{|S|^2 |E|^2} \sum_{j,l} (E_j - E_l)^{-1}$

For typical eigenvalue distribution goes with  $O(1/\log(|E|))$

# Fast Equilibration

Calculation only involves **4th** moments:

$$\frac{1}{T} \int_{U(|S||E|)} \text{tr} \left( \left( \text{tr}_E \left( U e^{-itD} U^\dagger \rho_0 U e^{itD} U^\dagger \right) - \rho_{\text{equi}} \right)^2 \right) \mu_{\text{Haar}}(dU)$$

Can replace Haar measure by an approximate unitary **4**-design

**Cor** For most Hamiltonians of the form  $UDU^\dagger$  with  $U$  a random quantum circuit of  $\mathbf{O}(n^2)$  size, small subsystems equilibrate fast.

# Open Questions

- What happens in higher dimensions?
- Fast scrambling conjecture (Hayden et al '11)  
Do  $O(\log(n))$ -depth random circuits equilibrate?  
(Brown, Fawzi '13) true for depth  $O(\log^2(n))$
- Equilibration for time-independent *local* Hamiltonians?  
(B., Ciwiklinski et al '11, Masanes et al '11, Vinayak, Znidaric '11) time-independent non-local Ham.)

# Open Questions

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- Fast scrambling conjecture (Hayden et al '11)  
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Thanks!